## Module 2 - Assembly

## Lecture 10: Genomics

Bioinformatics Algorithms CSC4181/6802

Most slides used are from Ben Langmead's Teaching Materials (www.langmead-lab.org/teaching-materials)

## Sequencing Technology

## First generation



[^0]
## Sanger Sequencing

$\square$


PCR with fluorescent, chain-terminating ddNTPs

Original
DNA sequence,
PCR amplified \& denatured

Size separation by capillary gel electrophoresis

Laser excitation \& detection by sequencing machine
https://www.sigmaaldrich.com/CA/en/technical-documents/proto col/genomics/sequencing/sanger-sequencing

## Sequencing Technology

## First generation



# Sanger sequencing <br> Maxam and Gilbert <br> Sanger chain termination 

Infer nucleotide identity using dNTPs, then visualize with electrophoresis
$500-1,000 \mathrm{bp}$ fragments

## Sequencing Technology


https://www.pacb.com/blog/the-evolution-of-dna-sequencing-tools/

## Sequencing by Synthesis



## Sequencing Technology


https://www.pacb.com/blog/the-evolution-of-dna-sequencing-tools/

## Sequencing Technology


https://www.pacb.com/blog/the-evolution-of-dna-sequencing-tools/

## PacBio Sequencing


https://www.pacb.com/wp-content/uploads/SMRT-Sequencing-Brochure-Deliverin g-highly-accurate-long-reads-to-drive-discovery-in-life-science.pdf

## Nanopore Sequencing


https://www.nature.com/articles/s41 587-021-01108-x/figures/1

## Sequencing Technology


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## Capturing measurement error: FASTQ



| Phred Quality Score | Probability of incorrect <br> base call | Base call accuracy |
| :--- | :--- | :--- |
| 10 | 1 in 10 | $90 \%$ |
| 20 | 1 in 100 | $99 \%$ |
| 30 | 1 in 1000 | $99.9 \%$ |
| 40 | 1 in 10000 | $99.99 \%$ |
| 50 | 1 in 100000 | $99.999 \%$ |

https://www.drive5.com/usearch/manual/fastq_files.html https://learn.gencore.bio.nyu.edu/ngs-file-formats/quality-scores/

## Assembly



Referenco genoma


How do we assemble puzzle without the benefit of knowing what the finished product should look like?
(That's what the Human Genome Project had to do!)

## De novo shotgun assembly



## Assembly

Whole-genome "shotgun" sequencing first copies the input DNA:

Input: GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT
Copy: GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT

Then fragments it:

| Fragment: | GGCGTCTA TATCTCGG | CTCTAGGCCCTC ATTTTTT |  |
| ---: | :--- | ---: | :--- |
|  | GGC GTCTATAT | CTCGGCTCTAGGCCCTCATTTTTT |  |
|  | GGCGTC TATATCT | CGGCTCTAGGCCCT | CATTTTTT |
|  | GGCGTCTAT ATCTCGGCTCTAG | GCCCTCA | TTTTTTT |

"Shotgun" refers to the random fragmentation of the whole genome; like it was fired from a shotgun

## Assembly

CTAGGCCCTCAATTTTT
CTCTAGGCCCTCAATTTTT
GGCTCTAGGCCCTCATTTTTT $|$

## Assembly

|  | CTAGGCCCTCAATTTTT |
| :--- | :--- |
|  | GGCGTCTATATCT |
|  | CTCTAGGCCCTCAATTTTT |
| Reconstruct this | TCTATATCTCGGCTCTAGG |
|  | GGCTCTAGGCCCTCATTTTT |
|  | CTCGGCTCTAGCCCCTCATTTT |
|  | TATCTCGACTCTAGGCCCTCA |
|  | GGCGTCGATATCT |
|  | TATCTCGACTCTAGGCC |
| GGCGTCTATATCTCG |  |
| $\longrightarrow ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ? ~$ |  |

Coverage

## CTAGGCCCTCAATTTTT

 CTCTAGGCCCTCAATTTTT GGCTCTAGGCCCTCATTTTTT CTCGGCTCTAGCCCCTCATTTT TATCTCGACTCTAGGCCCTCA TATCTCGACTCTAGGCCTCTATATCTCGGCTCTAGG
GGCGTCTATATCTCG
GGCGTCGATATCT
GGCGTCTATATCT
GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT
Coverage $=5$

Coverage
CTAGGCCCTCAATTTTT CTCTAGGCCCTCAATTTTT GGCTCTAGGCCCTCATTTTTT CTCGGCTCTAGCCCCTCATTTT TATCTCGACTCTAGGCCCTCA TATCTCGACTCTAGGCC
TCTATATCTCGGCTCTAGG
GGCGTCTATATCTCG
GGCGTCGATATCT
GGCGTCTATATCT
GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT
Coverage $=5$

## CTAGGCCCTCAATTTTT CTCTAGGCCCTCAATTTTT GGCTCTAGGCCCTCATTTTTT CTCGGCTCTAGCCCCTCATTTT TATCTCGACTCTAGGCCCTCA TATCTCGACTCTAGGCC TCTATATCTCGGCTCTAGG

GGCGTCTATATCTCG
GGCGTCGATATCT
GGCGTCTATATCT 35 bases
GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT
Average coverage $=177$ / $35 \approx 5$-fold

## TCTATATCTCGGCTCTAGG

## TATCTCGACTCTAGGCC

## TCTATATCTCGGCTCTAGG ||l|||| |||||| TATCTCGACTCTAGGCC

## First law of assembly

If a suffix of read $A$ is similar to a prefix of read $B$...

...then A and B might overlap in the genome

> TCTATATCTCGGCTCTAGG
> GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT TATCTCGACTCTAGGCC

## TCTATATCTCGGCTCTAGG |||||||||||| TATCTCGACTCTAGGCC $\uparrow$

Why the differences?

1. Sequencing errors
2. Ploidy: e.g. humans have 2 copies of each chromosome, and copies can differ


## Second law of assembly

More coverage leads to more and longer overlaps

```
CTAGGCCCTCAATTTTT CTCGGCTCTAGCCCCTCATTTT TCTATATCTCGGCTCTAGG GGCGTCGATATCT less coverage GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT CTAGGCCCTCAATTTTT GGCTCTAGGCCCTCATTTTTT CTCGGCTCTAGCCCCTCATTTT TATCTCGACTCTAGGCCCTCA
TCTATATCTCGGCTCTAGG
GGCGTCTATATCTCG
GGCGTCTATATCT

\section*{TCTATATCTCGGCTCTAGG ||||||| |||||| TATCTCGACTCTAGGCC}

\section*{TCTATATCTCGGCTCTAGG \(\|\|\|\|\|\|\|\|\) TATCTCGACTCTAGGCC}

\section*{TATCTCGACTCTAGGCC |||| |||||| || CTCGGCTCTAGCCCCTCAT}

\section*{Directed graph}


Directed graph


\section*{Overlap graph}

Each node is a read

\section*{CTCGGCTCTAGCCCCTCATTTT}

Draw edge \(A\)-> \(B\) when suffix of \(A\) overlaps prefix of \(B\)

\section*{Overlap graph}

Nodes: all 6-mers from GTACGTACGAT
Edges: overlaps of length \(\geq 4\)


\section*{Overlap graph}

Nodes: all 6-mers from GTACGTACGAT
Edges: overlaps of length \(\geq 4\)


\section*{Overlap Layout Consensus}


\section*{Finding overlaps}

Overlap: Suffix of \(X\) of length \(\geq l\) matches prefix of \(Y ; l\) is given
Naive: look in \(X\) for occurrences of \(Y\) 's length- \(l\) prefix. Extend matches to the right to confirm whether entire suffix of \(X\) matches.


See suffixPrefixMatch function in HW5 Q4 (Assembly Challenge)

\section*{Finding overlaps}

\section*{With suffix tree?}

Given a collection of strings \(S\), for each string \(x\) in \(S\) find all overlaps involving a prefix of \(x\) and a suffix of another string \(y\)

\section*{Finding overlaps with suffix tree}

Generalized suffix tree for \(\{\) "GACATA", "ATAGAC" \(\} \quad\) GACATA\$ \({ }_{0} A T A G A C \$ 1\)


\section*{Finding overlaps with suffix tree}

Generalized suffix tree for \(\{\) "GACATA", "ATAGAC" \(\} \quad\) GACATA\$ \({ }_{0} A T A G A C \$ 1\)


\section*{Finding overlaps with suffix tree}

Generalized suffix tree for \(\{\) "GACATA", "ATAGAC" \(\} \quad\) GACATA\$0ATAGAC\$1


\section*{Finding overlaps with suffix tree}


Assume for given string pair we report only the longest suffix/prefix match
Time to build generalized suffix tree: \(O(N)\)
... to walk down red paths: \(\quad O(N)\)
... to find \& report overlaps (green): O(a)
Overall:
\(\mathrm{O}(N+a)\)

\section*{Finding overlaps}

What about approximate suffix/prefix matches?

Dynamic programming

\section*{Finding overlaps with dynamic programming}

\section*{X: CTCGGCCCTAGG ||| ||||| \\ Y: GGCTCTAGGCCC}

Use global alignment recurrence and score function


How do we force it to find prefix / suffix matches?

\section*{Finding overlaps with dynamic programming}
\begin{tabular}{|c|c|c|c|c|c|}
\hline \(s(a, b)\) & A & C & G & T & - \\
\hline A & 0 & 4 & 2 & 4 & 8 \\
\hline C & 4 & 0 & 4 & 2 & 8 \\
\hline G & 2 & 4 & 0 & 4 & 8 \\
\hline T & 4 & 2 & 4 & 0 & 8 \\
\hline - & 8 & 8 & 8 & 8 & \\
\hline
\end{tabular}

How to initialize first row \& column so suffix of \(X\) aligns to prefix of \(Y\) ?

First column gets \(0 s\) (any suffix of \(X\) is possible)

First row gets \(\infty\) s (must be a prefix of \(Y\) )

Backtrace from last row
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline & & - & G & G & C & T & C & T & A & G & G & C & C & C \\
\hline & & 0 & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) & \(\infty\) \\
\hline & C & 0 & 4 & 12 & 20 & \multicolumn{9}{|l|}{\multirow[t]{4}{*}{STGGCCCTAGG
\(Y: \quad| || || |\)
\(Y: \quad\) GGCTCTAGGCCC}} \\
\hline & T & 0 & 4 & 8 & 14 & & & & & & & & & \\
\hline & C & & 4 & 8 & 8 & & & & & & & & & \\
\hline & G & 0 & & 4 & 12 & & & & & & & & & \\
\hline & G & 0 & 0 & & 8 & 10 & & <4 & <0 & 30 & So & 44 & K & \\
\hline \multirow[t]{7}{*}{\(X\)} & C & 0 & 4 & 4 & 9 & 8 & 16 & 18 & 26 & 30 & 34 & 36 & 44 & 52 \\
\hline & C & 0 & 4 & 8 & 4 & 2 & 8 & 16 & 22 & 30 & 34 & 34 & 36 & 44 \\
\hline & C & 0 & 4 & 8 & 8 & 6 & 2 & 10 & 18 & 26 & 34 & 34 & 34 & 36 \\
\hline & T & 0 & 4 & 8 & 10 & 8 & 8 & 2 & 10 & 18 & 26 & 34 & 36 & 36 \\
\hline & A & 0 & 2 & 6 & 12 & 14 & 12 & 10 & 2 & 10 & 18 & 26 & 34 & 40 \\
\hline & G & 0 & 0 & 2 & 10 & 16 & 18 & 16 & 10 & ? & 10 & 18 & 26 & 34 \\
\hline & G & 0 & 0 & 0 & 6 & 14 & 20 & 22 & 18 & 10 & ) & 10 & 18 & 26 \\
\hline
\end{tabular}

\section*{Finding overlaps with dynamic programming}

Say there are \(d\) reads of length \(n\), total length \(N=d n\), and \(a\) is total number of pairs with an overlap
\begin{tabular}{ll} 
\# overlaps to try: & \(O\left(d^{2}\right)\) \\
Size of each DP matrix: & \(O\left(n^{2}\right)\) \\
Overall: & \(O\left(d^{2} n^{2}\right)\), or \(O\left(N^{2}\right)\)
\end{tabular}

Contrast \(\mathrm{O}\left(N^{2}\right)\) with suffix tree: \(\mathrm{O}(N+a)\), but where \(a\) is worst-case \(\mathrm{O}\left(d^{2}\right)\)
Real-world overlappers mix the two; index filters out vast majority of non-overlapping pairs, dynamic programming used for remaining pairs

\section*{Overlap Layout Consensus}


\section*{Layout}

Overlap graph is big and messy. Contigs don't"pop out" at us.
Below: part of the overlap graph for
to_every_thing_turn_turn_turn_there_is_a_season
\(l=4, k=7\)


\section*{Layout}

Anything redundant about this part of the overlap graph?

Some edges can be inferred (transitively) from other edges
E.g. green edge can be inferred from blue


\section*{Layout}

Remove transitively inferrable edges, starting with edges that skip one node:


Before:


\section*{Layout}

Remove transitively inferrable edges, starting with edges that skip one node:


After:


\section*{Layout}

Now remove edges that skip one or two nodes:


After:


Even simpler

\section*{Layout}

Emit contigs corresponding to the non-branching stretches


\section*{Layout}

Must handle subgraphs that are spurious, e.g. because of sequencing error


Mismatch could be due to sequencing error or repeat. Since the path through \(\mathbf{b}\) ends abruptly we might conclude it's an error and prune \(\mathbf{b}\).

\section*{Overlap Layout Consensus}


\section*{Consensus}
\begin{tabular}{lll} 
TAGATTACACAGATTACTGA TTGATGGCGTAA CTA \\
TAGATTACACAGATTACTGACTTGATGGCGTAAACTA \\
TAG TTACACAGATTATTGACTTCATGGCGTAA CTA & Take reads that make \\
TAGATTACACAGATTACTGACTTGATGGCGTAA CTA & up a contig and line \\
them up \\
TAGATTACACAGATTACTGACTTGATGGCGTAA CTA
\end{tabular}

Complications: (a) sequencing error, (b) ploidy

\section*{Overlap Layout Consensus}


OLC drawbacks
Building overlap graph is slow. We saw \(\mathrm{O}(N+a)\) and \(\mathrm{O}\left(N^{2}\right)\) approaches.
Overlap graph is big; one node per read, \# edges can grow superlinearly with \# reads

Sequencing datasets are \(\sim\) 100s of millions or billions of reads```


[^0]:    Sanger sequencing
    Maxam and Gilbert
    Sanger chain termination

